Efficient arithmetic on (hyper-)elliptic curves over finite fields

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Overview

1. Arithmetic on hyperelliptic curves

2. Speed-up via endomorphisms

   (a) Koblitz curves

   (b) Generalized GLV

   (c) Trace zero subvariety

Group G suitable

1. scalar multiplication is fast,
   $\Rightarrow$ key-exchange is fast

2. group order can be computed efficiently,
   (avoid Pohlig-Hellman, . . .)

3. discrete logarithm problem is computationally hard,

4. representation is easy and compact.

Examples (group)

1. Diffie, Hellman (1976): $G = \mathbf{F}_p^*$, i.e.

$$G = \{1, 2, \ldots, p-1\}$$

2. $G = \mathbf{F}_q^*$, $q = p^l$

3. Miller/Koblitz (1985): $G = E(\mathbf{F}_q)$, group of points on an elliptic curve

4. Koblitz (1989): $G = \mathbf{Cl}(C/\mathbf{F}_q)$, ideal class group of a hyper-elliptic curve C over $\mathbf{F}_q$

Hyperelliptic curve

$$C : y^2 + h(x)y = f(x)$$

$h(x), f(x) \in \mathbf{F}_q[x]$, q prime power,
f monic, $\deg f = 2g + 1, \deg h \leq g$

nonsingular, i. e. not both partial derivatives zero for $(a, b) \in C/\bar{\mathbf{F}}_q$

C hyperelliptic curve of genus g

with at least one $\mathbf{F}_q$-rational Weierstraß point

(elliptic curves are curves of genus 1)

Group

Points do **not** form a group

use **divisors**, i. e. sums of points, of degree zero and reduce modulo principal divisors $\Rightarrow$ **divisor class group**.

$$\mathrm{Pic}^0(C/\mathbf{F}_q) = Div^0(\mathbf{F}_q)/P(\mathbf{F}_q)$$

class represented by degree zero divisor with $m \leq g$

$$\sum_{\substack{i=1 \\ P_i \in C \setminus \{\infty\}}}^{m} P_i - m\infty$$

Background

Divisor class group isomorphic to
ideal class group $\mathrm{Cl}(C/\mathbf{F}_q)$ of $\mathbf{F}_q(x,y)/(C)$.

Elements D represented by two polynomials

$$D = [u(x), v(x)]; u, v \in \mathbf{F}_q[x],$$

$\deg v < \deg u \leq g, \quad u | v^2 + vh - f.$

$\Rightarrow$ compact representation

isomorphism via

$P_i = (x_i, y_i) \Leftrightarrow u(x_i) = 0, v(x_i) = y_i$ with multiplicity

isomorphic to Jacobian – g-dimensional Abelian variety over $\mathbf{F}_q$

Arithmetic on hyperelliptic curves (Cantor/Koblitz)

Composition

IN: $D_1 = [u_1, v_1], D_2 = [u_2, v_2], C : y^2 + h(x)y = f(x)$
OUT: $D = [u, v]$ semireduced with $D \sim D_1 + D_2$

1. compute $d_1 = \gcd(u_1, u_2) = e_1 u_1 + e_2 u_2$

2. compute $d = \gcd(d_1, v_1 + v_2 + h) = c_1 d_1 + c_2(v_1 + v_2 + h)$

3. let $s_1 = c_1 e_1, s_2 = c_1 e_2, s_3 = c_2$
   /* i.e. $d = s_1 u_1 + s_2 u_2 + s_3(v_1 + v_2 + h)$

4. $u = \frac{u_1 u_2}{d^2} \quad v = \frac{s_1 u_1 v_2 + s_2 u_2 v_1 + s_3(v_1 v_2 + f)}{d} \bmod u$

Reduction

`IN:` $D = [u, v]$ semireduced

`OUT:` $D' = [u', v']$ reduced with $D \sim D'$

1. let $u' = \frac{f - vh - v^2}{u}$
   $b' = (-h - v) \bmod u'$

2. if $\deg u' > g$ put $u := u', v := v'$ goto step 1

3. make u monic

Arithmetic in $\mathbf{Cl}(C/\mathbf{F}_q)$

do better with explicit formulae – genus 2

2000 Harley (odd char.)
2001 L. (arbitrary char.)
2001 Matsuo, Chao, Tsujii (faster)

$\Rightarrow$ 2 inversions

2002 Miyamoto, Doi, Matsuo, Chao, Tsujii
2002 Takahashi
2002 L. (arbitrary char.)
2002 Sugizaki, Matsuo, Chao, Tsujii (even char.)

$\Rightarrow$ 1 inversion

genus 3

2002 Kuroki, Gonda, Matsuo, Chao, Tsujii
2002 Pelzl, Guyot & Patankar

$\Rightarrow$ 1 inversion

Different Coordinates (genus 2)

affine $\mathcal{A}$:

$$D \sim [x^2 + u_1 x + u_0, v_1 x + v_0] \sim [u_1, u_0, v_1, v_0]$$

projective $\mathcal{P}$: Miyamoto, Doi, Matsuo, Chao, and Tsujii; L.; all 2002

$$D \sim [U_1, U_0, V_1, V_0, Z] \sim [U_1/Z, U_0/Z, V_1/Z, V_0/Z]$$

new $\mathcal{N}$: L. 2002, contains also mixed coordinates

$$D \sim [U_1, U_1, V_1, V_0, Z_1, Z_2] \sim [U_1/Z_1^2, U_0/Z_1^2, V_1/(Z_1^3 Z_2), V_0/(Z_1^3 Z_2)]$$

$\mathcal{P}$ and $\mathcal{N}$: **no inversions**

Comparison - odd characteristic

Doubling		Addition	
operation	costs	operation	costs
$2\mathcal{N} = \mathcal{P}$	7S, 38M	$\mathcal{N} + \mathcal{N} = \mathcal{P}$	7S, 51M
$2\mathcal{P} = \mathcal{P}$	6S, 38M	$\mathcal{N} + \mathcal{P} = \mathcal{P}$	4S, 51M
$2\mathcal{N} = \mathcal{N}$	7S, 34M	$\mathcal{N} + \mathcal{N} = \mathcal{N}$	7S, 47M
$2\mathcal{P} = \mathcal{N}$	6S, 34M	$\mathcal{N} + \mathcal{P} = \mathcal{N}$	5S, 45M
$2\mathcal{A} = \mathcal{P}$	5S, 25M	$\mathcal{P} + \mathcal{P} = \mathcal{P}$	5S, 45M
$2\mathcal{A} = \mathcal{N}$	5S, 21M	$\mathcal{P} + \mathcal{P} = \mathcal{N}$	5S, 41M
		$\mathcal{A} + \mathcal{N} = \mathcal{P}$	5S, 41M
		$\mathcal{A} + \mathcal{P} = \mathcal{P}$	3S, 39M
		$\mathcal{A} + \mathcal{N} = \mathcal{N}$	5S, 36M
		$\mathcal{A} + \mathcal{P} = \mathcal{N}$	3S, 35M
		$\mathcal{A} + \mathcal{A} = \mathcal{N}$	3S, 25M
$2\mathcal{A} = \mathcal{A}$	1I, 5S, 22M	$\mathcal{A} + \mathcal{A} = \mathcal{A}$	1I, 3S, 22M

Implementations and Comparison

software: for $\log_2 p \sim 40 - 60$ inversions cheap,
use $\mathcal{A}$ for $g = 1, 2$. Comparison depends on I/M genus
Upcoming paper with R. Avanzi (Essen), library for $p \sim 160$ bit.

hardware: usually expensive inversions,
compute multiple of affine point via $2\mathcal{N} = \mathcal{N}, \mathcal{A} + \mathcal{N} = \mathcal{N}$,
compute multiple of non-normalized point via
$2\mathcal{N} = \mathcal{N}$, $\mathcal{N} + \mathcal{P} = \mathcal{N}$.

Implementation by K. Nguyen (Philips) presented at ECC 2002
using $\mathcal{P}$ is half as fast as elliptic – but same order . . . and elliptic curves are really optimized there.

Koblitz curves (subfield curves)

$$C : y^2 + h(x)y = f(x)$$

hyperelliptic curve with
$h(x), f(x) \in \mathbf{F}_q[x]$, q **small**
consider curve over $\mathbf{F}_{q^n}$

Example:

$C : y^2 + xy = x^5 + x^2 + 1$

hyperelliptic curve of genus 2 over $\mathbf{F}_2$
hyperelliptic curve of genus 2 over $\mathbf{F}_{2^n}$

History: Koblitz '89; Günther, L., Stein '00; L. 01 (Ph.D.)

Computing m-folds

Classical: **binary (double-and-add)**

Density:

$$\sim 1/2$$

Length of expansion of $m \in \mathbf{Z}$:

$$\lfloor \log_2 m \rfloor$$

for $m \leq |\mathbf{Cl}(C/\mathbf{F}_q)| \sim q^{ng}$:

$$\textbf{Complexity} \sim \tfrac{3}{2} \log_2(q^{gn})$$

For $q = 2$: $\sim \frac{3}{2} gn$

Frobenius endomorphism σ :

$P = (a, b) \in C(\mathbf{F}_{q^n})$

$$\Rightarrow \sigma(P) := (a^q, b^q) \in C(\mathbf{F}_{q^n})$$

Normal basis

Basis of $\mathbf{F}_{q^n} = \{\theta, \theta^q, \theta^{q^2}, \ldots, \theta^{q^{n-1}}\}$

Element $d \in \mathbf{F}_{q^n}$ represented by

$$d = (d_0, d_1, \ldots, d_{n-1}), \quad d_i \in \mathbf{F}_q$$

$$d^q = (d_{n-1}, d_0, d_1, \ldots, d_{n-2})$$

Remark: d^q realized by **cyclic shifting**,
q-th power also fast in polynomial basis

Frobenius endomorphism on $\mathrm{Cl}(C/\mathbf{F}_q)$

Theorem 1 $D = [u(x), v(x)]$

$$\sigma(D) = [\sigma(u(x)), \sigma(v(x))],$$

where $\sigma(\sum d_i x^i) = \sum d_i^q x^i$.

Execution of Frobenius endomorphism takes at most $2g$ cyclic shiftings

Example

C curve of genus 2 defined over $\mathbf{F}_2$

$$D = [x^2 + u_1 x + u_0, v_1 x + v_0] \quad \sigma(D) = [x^2 + u_1^2 x + u_0^2, v_1^2 x + v_0^2]$$

Characteristic polynomial of Frobenius endomorphism

$$P(T) = T^{2g} + a_1 T^{2g-1} + \cdots + a_g T^g + \cdots + a_1 q^{g-1} T + q^g$$

integer coefficients, degree $2g$

Examples for genus 2 over $\mathbf{F}_2$

Equation of C	$P(T)$
$v^2 + v = u^5 + u^3$	$T^4 + 2T^3 + 2T^2 + 4T + 4$
$v^2 + v = u^5 + u^3 + 1$	$T^4 - 2T^3 + 2T^2 - 4T + 4$
$v^2 + v = u^5 + u^3 + u$	$T^4 + 2T^2 + 4$
$v^2 + uv = u^5 + 1$	$T^4 + T^3 + 2T + 4$
$v^2 + uv = u^5 + u^2 + 1$	$T^4 - T^3 - 2T + 4$
$v^2 + (u^2 + u + 1)v = u^5 + u^4 + u^3$	$T^4 + T^2 + 4$
$v^2 + (u^2 + u)v = u^5 + u^4 + u$	$T^4 - T^2 + 4$
$v^2 + (u^2 + u + 1)v = u^5 + u^4$	$T^4 + 2T^3 + 3T^2 + 4T + 4$
$v^2 + (u^2 + u + 1)v = u^5 + u^4 + 1$	$T^4 - 2T^3 + 3T^2 - 4T + 4$

$P(T)$ **easy** to compute for **Koblitz curves**

Consequences

$$-q^g D = \sigma^{2g}(D) + a_1\sigma^{2g-1}(D) + \cdots + a_1 q^{g-1}\sigma(D)$$

for $D \in \mathbf{Cl}(C/\mathbf{F}_q)$.

To compute $mD = \sum u_i \sigma^i(D)$ use

τ-adic expansion

$$m = \sum u_i \tau^i,$$

where τ complex root of $P(T)$,

$u_i \in R$: **set of coefficients**.

τ-adic

input: $D \in \mathbf{Cl}(C/\mathbf{F}_{q^n})$, $m = \sum_{i=0}^{l-1} u_i \tau^i, u_i \in R$.

output: $H = mD$

1. precompute rD for $r \in R,\ r > 0$.

2. $H \leftarrow u_{l-1} D$

3. For $i = l-2$ to 0 do
   $H \leftarrow \sigma(H)$
   if $(u_i \neq 0)$ $H \leftarrow H + u_i D$

4. return(H)

τ-adic expansion

Lemma 1 $c = c_1 + c_2\tau + \cdots + c_{2g}\tau^{2g-1} \in \mathbf{Z}[\tau]$

$$\tau | c \Leftrightarrow q^g | c_1$$

Necessary for expansion:

$R \supseteq \{0, \pm 1, \ldots, \pm\lceil \frac{q^g-1}{2} \rceil\}$

Note:

Taking negative is for free.

$$D = [u(x), v(x)] \Rightarrow -D = [u(x), -h(x) - v(x)]$$

Comparison

Operations in **binary** representation:

$$\sim \frac{3}{2}\, g\ n\ \log_2 q \geq \frac{3}{2}\, g\ n$$

Operations in τ-**adic** representation:

- $(q^g - 2)/2$ precomputations

$$\sim \frac{q^g - 1}{q^g}\ n < n$$

- $(q^g - 1)q^g/2$ precomputations

$$\sim \frac{q^g - 1}{2q^g - 1}\ n < \frac{1}{2}\ n$$

Comparison
(same number of precomputations as signed binary)

$q = 2$:

g	binary window	τ-adic $w = 1$	speed-up	binary window	τ-adic $w = 2$	speed-up
2	$11/4n$	$3/4n$	11/3	$31/12n$	$3/7n$	$217/36 \sim 6$
3	$31/8n$	$7/8n$	31/7	$573/160n$	$7/15n$	$1719/224 \sim 7.6$
4	$79/16n$	$15/16n$	79/15	$1023/224n$	$15/31n$	$10571/1120 \sim 9.4$

$q = 5$: (friendly)

g	w_{bin}	binary window	τ-adic $w = 1$	speed-up	w_{bin}	binary window	τ-adic $w = 2$	speed-up
2	4	$47/8n$	$24/25n$	6	9	$6n$	$24/49n$	12
3	7	$511/64n$	n	8	13	$8n$	$124/249n$	16
4	9	$10n$	n	10	18	$10n$	$1/2n$	20

Other endomorphisms – generalization of GLV
Gallant, Lambert, and Vanstone

(Park, Jeong, Kim, and Lim; Sica, Ciet, and Quisquater; 2002)

ϕ endomorphism of C, $P_\phi(T)$ characteristic polynomial of ϕ,

$d = \deg P \leq 2g$, l group order.
split

$$m = m_0 + m_1\phi + \cdots + m_{d-1}\phi^{d-1}, |m_i| \sim l^{1/d}$$

for $d = 2$ use joint-sparse form (JSF) of exponent,
otherwise interleaved multiplications.

Both significantly faster than binary method

Combination of GLV and expansion – so far $g = 1$

(Ciet, L., Sica, Quisquater)

ϕ satisfies

$$T^2 + rT + s$$

use GLV

$$m = m_0 + m_1\phi, \quad |m_i| \sim l^{1/2}$$

expand

$$m_0 + m_1\phi = k_0 + k_1\phi + k_2\phi^2 + \cdots, \ k_i \in \mathcal{R}, \ \text{length} \sim \log_s m$$

precompute $r_i P$ for $r_i \in \mathcal{R} = \{0, \pm 1, \ldots, \lceil (s-1)/2 \rceil\}$

together with ϕ-JSF $\Rightarrow$ further speed-up

Trace-zero subvariety G
('Koblitz curves' for small extension)

genus 1 or 2 curve defined over $\mathbf{F}_p$, considered over $\mathbf{F}_{p^3}$

trace zero property

$$D \in G \Leftrightarrow D \in \mathbf{Cl}(C/\mathbf{F}_{p^3}) \text{ and } D + \sigma(D) + \sigma^2(D) = 0,$$

subgroup of $\mathbf{Cl}(C/\mathbf{F}_{p^3})$

History: Frey '98, Naumann '99 ($g = 1$),
L. '01 ($g = 2$), Blady '02 ($g = 1$)
Silverberg/Rubin '02 (for ID-based systems)

Background

dimension $2g$, group size from $P(T)$,

$$\#G \sim \#\mathbf{Cl}(C/\mathbf{F}_{p^3})/\#\mathbf{Cl}(C/\mathbf{F}_p) \sim p^{2g}$$

Theorem 2 *Let* $2, 3 \not| \, l \,|\, \#\mathbf{Cl}(C/\mathbf{F}_{p^3})$. *Then the nontrivial elements in the subgroup of order* l *of the trace zero variety* G *are the divisor classes represented by*

$$P_1 + P_2 - 2\infty,$$

where $P_1 \neq P_2, \sigma(P_2), \sigma^2(P_2)$ *and not both* $P_1, P_2 \in C(\mathbf{F}_p)$.

Scalar multiples in G

instead of computing mD for $0 \leq m < l,\ l \sim p^{2g}$ group oder $\langle D \rangle \subseteq G$ use

$$r_0 D + r_1 \sigma(D),\ |r_i| \sim \sqrt{m}.$$

and Joint Sparse Form (JSF) of r_0, r_1.

Avoid collisions ($g = 2$)

$$r_0, r_1 < \mathbf{r} := \min \left\{ \left\lfloor \frac{l}{m} \right\rfloor, \frac{p^2 - a_2 + a_1}{\gcd(p^2 - a_2 + a_1, a_1 p - a_2 + 1)} \right\}.$$

where $m = \max\{p^2 + a_1 p - 2a_2 + a_1 + 1, p^2 + a_1 - a_1 p - 1\}$.

The $\mathbf{r}^2$ classes $r_0 D + r_1 \sigma(D),\ 0 \leq r_i < \mathbf{r}$ are distinct.

Comparison

	Elliptic, in $\mathbf{F}_{p''}$			Genus 2, in $\mathbf{F}_{p'}$			G ($g = 2$), in $\mathbf{F}_p$		
	Inv.	Sqr.	Mult.	Inv.	Sqr.	Mult.	Inv.	Sqr.	Mult.
Add.	1	1	2	1	3	22	1	21	141
Doub.	1	2	2	1	5	22	1	33	141
m-fold	6λ	10λ	12λ	6λ	26λ	132λ	3.5λ	97.5λ	501.5λ
NAF	4.8λ	8.8λ	9.6λ	4.8λ	22.4λ	105.6λ	3λ	87λ	423λ

using efficient arithmetic in $\mathbf{F}_{p^3}$

(NAF includes same number of precomputations)

The END